

[DOT<sup>+</sup>23]

# Doing Math with Computers for Fun and Profit

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# Hi, I'm Ryan!

- Junior in CS, on a mission to take all the CS theory courses
- I also do web dev :(
- Interests: randomized and streaming algorithms, (some) systems
- Will be presenting later this semester!



# Section 1

REU



## What is an REU?

- Research Experience for Undergraduates
- Get paid to do research in Math, CS, Engineering, Science, etc., over the summer
- See how other schools do things, meet new people
- Maybe even get a paper out of it!



## How do I find them?

Save these slides for later!

- NSF List
- Math Programs
- This spreadsheet



## How do I apply to them?

- Personal Statement
- 2 Letters of Rec
- Resume / CV
- Most deadlines are early – mid March
  - ▶ Start drafting in Winter Break



## Tips & Tricks

- Most are government funded which means usually US citizens get funding.
  - ▶ It is possible for non-US citizens to get funding in some special cases.
- Get your letter writers to read your personal statement.
- There are other options outside of REUs ([MSR](#), [TTIC](#), [EPFL](#), [ETH Zürich](#), [Max Planck](#), [SCAMP](#), Independent Study, etc).
- Most people don't apply to enough REUs.

Any other questions?



Questions?





## Section 2

### Introduction to Group Theory



# Groups and Group Actions

- A *group* is an object in the category of groups
- A *group action* is a functor from a 1-groupoid to the category of sets



## What is a Group?

Groups are one of the most ubiquitous objects in all of math. They generalize structures with some sort of addition/multiplication.

### Definition

A *group* is a set  $G$  with an operation  $\cdot : G \times G \rightarrow G$  such that

- $\cdot$  is associative:  $(x \cdot y) \cdot z = x \cdot (y \cdot z)$
- There exists an identity  $e$  such that  $g \cdot e = g = e \cdot g$
- Every element  $g$  has an inverse  $g^{-1}$  such that  $g \cdot g^{-1} = e = g^{-1} \cdot g$

We will usually just write  $x \cdot y$  as  $xy$



## Important Examples of Groups

Consider the set  $S_n$  of bijections  $\sigma: [n] \rightarrow [n]$  where  $[n] = \{1, \dots, n\}$ . This forms a group with “multiplication” using composition

- Composing bijections with each other yields a bijection
- Identity:  $\text{id}(i) = i$  for all  $1 \leq i \leq n$
- Inverses:  $\sigma$  has an inverse  $\sigma^{-1}$  such that  $\sigma \circ \sigma^{-1} = \text{id}$



## Important Examples of Groups

Recall that the integers mod  $p$  are  $\mathbb{Z}_p = \{1, 2, \dots, p\}$  with addition and multiplication done modulo  $p$ . Consider the set  $GL(n, p)$  of all  $n \times n$  matrices with entries in  $\mathbb{Z}_p$  with non-zero determinant. This forms a group with matrix multiplication

- Multiplying two matrices with non-zero determinant yields a matrix with non-zero determinant since  $\det(AB) = \det(A) \cdot \det(B)$
- Identity:  $I_n$  with 1s on the diagonal and 0s elsewhere
- Inverses: Matrices have an inverse if and only if they have non-zero determinant, so each  $A$  has an inverse  $A^{-1}$  such that  $A \times A^{-1} = I_n$

If you've taken Linear Algebra, these are just invertible linear transformations!



## Group Isomorphism

Consider the following two groups:

$$G = \left\{ \text{id} = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 2 & 3 \end{pmatrix}, \sigma = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 3 & 1 \end{pmatrix}, \sigma^2 = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 1 & 2 \end{pmatrix} \right\} \subseteq S_3$$

$$\mathbb{Z}_3 = \{ \quad \quad 0, \quad \quad \quad 1, \quad \quad \quad 2 \quad \quad \}$$

So  $G$  are some permutations with the operation of composition and  $\mathbb{Z}_3$  are integers modulo 3 where the operation is addition but we keep the remainder after division by 3. So  $2 + 2 \equiv 1 \pmod{3}$ .

**Question:** In what sense are these two groups the same group?

**Answer:** Mapping  $\text{id} \mapsto 0$ ,  $\sigma \mapsto 1$ ,  $\sigma^2 \mapsto 2$  preserve operations!

Notice that  $\sigma \circ \sigma = \sigma^2$  and  $1 + 1 \equiv 2 \pmod{3}$ .

Similarly,  $\sigma \circ \sigma \circ \sigma = \text{id}$  and  $1 + 1 + 1 \equiv 0 \pmod{3}$ .



## Group Actions

We want to study how a group  $G$  interacts with other sets. Let  $\Omega$  be some set.

### Definition

Then a *group action* is an operation  $\cdot : G \times \Omega \rightarrow \Omega$  such that

- $e \cdot x = x$ , for all  $x \in \Omega$
- $g \cdot (h \cdot x) = (gh) \cdot x$ , for all  $g, h \in G$ , and for all  $x \in \Omega$

We write  $G \curvearrowright \Omega$ .

To prevent confusion with the group operation in  $G$ , we will keep the  $\cdot$  when talking about actions.



## Important Examples of Group Actions

Let  $G = S_n$  and  $\Omega = [n]$ .

- Say  $\sigma: [n] \rightarrow [n] \in S_n$  and  $i \in [n]$ . What would be a good choice of action  $\sigma \cdot i$ ?
- $\sigma \cdot i := \sigma(i)$

Now let  $G$  be a group of invertible linear transformations from a vector space  $V \rightarrow V$ .

- Say  $T: V \rightarrow V \in G$  and  $v \in V$ . What would be a good choice of action  $T \cdot v$ ?
- $T \cdot v := T(v)$





## Orbits and Stabilizers

We want to study the structure of  $G \curvearrowright \Omega$ .

### Definition

The *orbit* of  $\alpha \in \Omega$  is the set  $G \cdot \alpha = \{ g \cdot \alpha \mid g \in G \}$

Every element of  $\Omega$  belongs in some orbit. It turns out the orbits partition  $\Omega$ .

### Definition

The *stabilizer* of  $\alpha \in \Omega$  is the set  $G_\alpha = \{ g \in G \mid g \cdot \alpha = \alpha \}$

**Exercise:** Stabilizers are subgroups of  $G$



## Example

$$G = \left\{ \text{id} = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 1 & 2 & 3 & 4 \end{pmatrix}, \sigma = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 2 & 3 & 1 & 4 \end{pmatrix}, \sigma^2 = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 3 & 1 & 2 & 4 \end{pmatrix} \right\} \subseteq S_4$$

This is the same group as before, but now 4 is a valid input and we don't do anything to it. **Exercise:** Check that this is a subgroup of  $S_4$ .

Consider  $G \curvearrowright [4]$ .

- $G \cdot 1 = \{1, 2, 3\}$
- $G \cdot 4 = \{4\}$
- $G_1 = \{\text{id}\}$
- $G_4 = G$



# Group Classification

The *Group Classification Problem* is the problem of identifying groups satisfying some property “up to” isomorphism.

- This is one of the hardest problems in all of group theory.
- Even checking if two finite groups are isomorphic is difficult for computer.
- The classification of the finite simple groups took tens of thousands of pages written by over 100 authors between 1955 and 2004.



## Rank

- Let  $G$  be some group of permutations in  $S_n$  and consider  $G \curvearrowright [n]$  such that  $G$  only has one orbit.
- Let  $G_0$  be the stabilizer of some element of  $[n]$ . It turns out it doesn't matter which one.

### Definition

The *rank* of  $G$  is the number of orbits of  $G_0 \curvearrowright [n]$ .

This is a sort of measurement of the “reach” of stabilizer subgroups of  $G$ .



## Section 3

# Groups, Algorithms, and Programming



## Structure

- Let  $G$  be some group of permutations in  $S_n$  and consider  $G \curvearrowright [n]$  such that  $G$  only has one orbit.
- Let  $G_0$  be the stabilizer of some element of  $\Omega$ .
- **Result 1:** This is the same as considering  $G_0 \curvearrowright V$  where  $G_0$  is now a group of linear transformations and  $V$  is a vector space  $\mathbb{F}_p^k$ .
- $\#$  Orbits of  $G_0 \curvearrowright [n] = \#$  Orbits of  $G_0 \curvearrowright \mathbb{F}_p^k$ .
- **Result 2:**  $G_0$  must contain a certain subgroup  $E$  of order  $q^{2m+1}$ .



## Making Change using Group Theory

So now we can consider  $G_0 \curvearrowright \mathbb{F}_p^k$  and  $E \subseteq G_0$  where  $|E| = q^{2m+1}$ . This gives us a nice set of parameters.

1. We find a value  $B(p, k, q, m)$  such that  $|G_0|$  divides  $B$
2. A theorem in group theory tells us that the size of orbits of  $G_0$  divides  $|G_0|$ , so they divide  $B$ 
  - ▶ Let  $d_1, \dots, d_t$  be the divisors of  $B$
3. We know there is one orbit of size 1 and the sizes of the other orbits must sum up to  $p^k - 1$

**Result 3:** We can get a lower bound on rank by solving the *Change Making Problem* with coins  $d_1, \dots, d_t$  and target value  $p^k - 1$ .



## Making Change using Group Theory

In the *Change-making Problem*, we are given coins from some set of denominations  $d_1, \dots, d_t$  and a target value  $T$ , we want to “make change” for  $T$  using as few coins as possible

- We have a fixed set of possible sizes of orbits and a target value  $p^k - 1$
- We know the orbits partition this target value
- A worst case lower bound is the most efficient packing as possible
- Thus we want to solve the Change Making Problem with coins  $d_1, \dots, d_t$  and target value  $p^k - 1$ .





## Inductively Making Change

Let  $coins = [d_1, \dots, d_n]$  be a sorted list of denominations of coins. Let  $NUMCOINS(i, c)$  be the smallest possible number coins of denomination  $[d_1, \dots, d_c]$  needed make change for  $i$

- If  $coins = [1, 3, 5, 7]$  then  $NUMCOINS(10, 1) = 10$  but  $NUMCOINS(10, 4) = 2$



## Inductively Making Change

Let  $coins = [d_1, \dots, d_n]$  be a sorted list of denominations of coins. Let  $NUMCOINS(i, c)$  be the smallest possible number coins of denomination  $[d_1, \dots, d_c]$  needed make change for  $i$ .

$NUMCOINS(i, c) =$

$$\begin{cases} \infty & c = 0 \\ NUMCOINS(i, c - 1) & i < coins[c] \\ 1 & i = coins[c] \\ \min \{ NUMCOINS(i, c - 1), 1 + NUMCOINS(i - coins[c], c) \} & \text{otherwise} \end{cases}$$



## A (very high level) Overview of the Whole Paper

1. Define the parameters  $p, k, q, m$
2. Do a bunch of pure math to get finite bounds on these parameters
3. Enumerate all possible sets of parameters and keep the ones that have a lower bound  $\leq 6$



## A (very high level) Overview of the Whole Paper

For each set of valid parameters  $p, k, q, m$ . Let  $N$  = the largest possible  $N$  such that  $N \curvearrowright \mathbb{F}_p^k$

1. Check if the subgroup  $E$  with  $|E| = q^{2m+1}$  is contained in  $N$   
(HARD!)
2. Check if  $N$  has rank  $\leq 6$
3. Enumerate all possible subgroups of  $N$  (HARD!)
4. Repeat for each subgroup



## How?

- All of this was done in a programming language called **GAP**: Groups, Algorithms, and Programming
- GAP is just one of many computational algebra systems
  - ▶ SageMath (Built on top of Python)
  - ▶ Mathematica
  - ▶ Magma (popular in Cryptography)
  - ▶ Macauley2 (Created at UIUC!)
- Hard computations were done on AWS.
- These techniques extend to higher ranks but computational resources are a large issue.



## More Details?

- Check out the paper (linked on my website [anakin-dey.com](https://anakin-dey.com))
- Come to the Undergraduate Math Seminar (details coming soon)
- Ask me in the Discord!



Questions?



*Algebra is the offer made by the devil to the mathematician...All you need to do, is give me your soul: give up geometry*

— Michael Atiyah ([1979](#))





# Bibliography



Anakin Dey, Kolton O'Neal, Duc Van Khanh Tran, Camron Upshur, and Yong Yang.

Classifying primitive solvable permutation groups of rank 5 and 6, 2023.

