Welcome to SIGma SIGma

## Outline

Officers in No Particular Order

Computing Fibonacci

## Anakin

- Math Major
- SIGPwny Crypto ${ }^{1}$ Gang + Admin team
- CA for CS $173+$ CS 475
- Research with Sam

[^0]
## Sam

- Summer Amazon Intern
- CS Major
- Doing CS 374 Course Dev
- Doing Theory Research with Sariel Har-Peled
- Research with Anakin


## Lou

- CS Major
- Current CS 225 CA (past CS 125 and 374 CA )
- Senior, selling soul to finance after this semester


## Aditya

- ECE/Math double degree.
- Quantum error correction research w/Prof. Milenkovic.
- CA for ECE 391.
- Other interests: FP, PL, Crypto.


## Hassam

- Intern at Amazon over the summer
- CS Major (takes math classes for fun ???)
- SIGPwny Crypto Gang + Admin team + Infra lead
- CA for CS 233, CS 173
- Compiler research


## Phil

- CS/Ling Major, Senior
- CA for CS 233
- SIGecom - game theory, economics, and computation

Section 2

Computing Fibonacci

Recursive

$$
F_{n}= \begin{cases}0 & n=0 \\ 1 & n=1 \\ F_{n-1}+F_{n-2} & n \geq 2\end{cases}
$$

## Recursive

$$
F_{n}= \begin{cases}0 & n=0 \\ 1 & n=1 \\ F_{n-1}+F_{n-2} & n \geq 2\end{cases}
$$

| $F_{0}$ | $F_{1}$ | $F_{2}$ | $F_{3}$ | $F_{4}$ | $F_{5}$ | $F_{6}$ | $F_{7}$ | $F_{8}$ | $F_{9}$ | $F_{10}$ | $F_{11}$ | $F_{12}$ | $F_{13}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 1 | 2 | 3 | 5 | 8 | 13 | 21 | 34 | 55 | 89 | 144 | 233 |

## Recursive Computation



Figure: From [Eri19]

## Can We Do Better?

We can use 12 multiplications to compute $x^{13}$ as follows:

$$
x \rightarrow x^{2} \rightarrow x^{3} \rightarrow x^{4} \rightarrow x^{5} \rightarrow x^{6} \rightarrow x^{7} \rightarrow x^{8} \rightarrow x^{9} \rightarrow x^{10} \rightarrow x^{11} \rightarrow x^{12} \rightarrow x^{13}
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But if we first compute powers as such

$$
\begin{aligned}
& x^{2} \leftarrow x \cdot x \\
& x^{4} \leftarrow x^{2} \cdot x^{2} \\
& x^{8} \leftarrow x^{4} \cdot x^{4}
\end{aligned}
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Using these we compute $x^{8} \cdot x^{4} \cdot x^{1}=x^{13}$ in just 5 total multiplications. We can generalize this using binary

| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ |
| :--- | :--- | :--- | :--- |
| 8 | 4 | 2 | 1 |

Building an Algorithm

$$
13=8+4+1=\mathbf{1 1 0 1}_{2}
$$

## Building an Algorithm

$$
13=8+4+1=\mathbf{1 1 0 1}_{2}
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| Step | Bit | Power | Result |
| :---: | :---: | :---: | :---: |
| 0 |  |  | 1 |

## Building an Algorithm

$$
13=8+4+1=\mathbf{1 1 0 1}_{2}
$$

| Step | Bit | Power | Result |
| :---: | :---: | :---: | :---: |
| 0 |  |  | 1 |
| 1 | $\mathbf{1}$ | $x$ | $x$ |

## Building an Algorithm

$$
13=8+4+1=\mathbf{1 1 0 1}_{2}
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| Step | Bit | Power | Result |
| :---: | :---: | :---: | :---: |
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| 1 | $\mathbf{1}$ | $x$ | $x$ |
| 2 | $\mathbf{0}$ | $x^{2}$ | $x$ |

## Building an Algorithm

$$
13=8+4+1=\mathbf{1 1 0 1}_{2}
$$

| Step | Bit | Power | Result |
| :---: | :---: | :---: | :---: |
| 0 |  |  | 1 |
| 1 | $\mathbf{1}$ | $x$ | $x$ |
| 2 | $\mathbf{0}$ | $x^{2}$ | $x$ |
| 3 | $\mathbf{1}$ | $x^{4}$ | $x^{5}$ |

## Building an Algorithm

$$
13=8+4+1=\mathbf{1 1 0 1}_{2}
$$

| Step | Bit | Power | Result |
| :---: | :---: | :---: | :---: |
| 0 |  |  | 1 |
| 1 | $\mathbf{1}$ | $x$ | $x$ |
| 2 | $\mathbf{0}$ | $x^{2}$ | $x$ |
| 3 | $\mathbf{1}$ | $x^{4}$ | $x^{5}$ |
| 4 | $\mathbf{1}$ | $x^{8}$ | $x^{13}$ |


|  | $\quad \operatorname{POWER}(x, n):$ |
| :--- | :---: |
| $1:$ | $\operatorname{curr} \leftarrow 1$ |
| $2:$ | for $i \leftarrow 1 \ldots n:$ |
| $3:$ | $\operatorname{curr} \leftarrow \operatorname{curr} * x$ |
| $4:$ | return curr |


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| $4:$ | return curr |


|  | SQUAREMULTPOWER $(\mathrm{x}, \mathrm{n}):$ |  |
| :---: | :---: | :---: |
| $1:$ | res $\leftarrow 1$ |  |
| $2:$ | power $\leftarrow x$ |  |
| $3:$ | for bit in BINARY $(\mathrm{n}):$ |  |
| $4:$ | if bit $=1:$ |  |
| $5:$ | res $\leftarrow$ res $*$ power |  |
| $6:$ | power $\leftarrow$ power $*$ power |  |
| $7:$ | return res |  |

## Matrices

We have the following two linear equations

$$
\begin{aligned}
F_{n} & =F_{n-1}+F_{n-2} \\
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We can represent this as follows using matrices

$$
\left[\begin{array}{c}
F_{n} \\
F_{n-1}
\end{array}\right]=\left[\begin{array}{ll}
1 & 1 \\
1 & 0
\end{array}\right]\left[\begin{array}{l}
F_{n-1} \\
F_{n-2}
\end{array}\right]
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F_{n-2}
\end{array}\right]=\left[\begin{array}{ll}
1 & 1 \\
1 & 0
\end{array}\right]^{2}\left[\begin{array}{l}
F_{n-2} \\
F_{n-3}
\end{array}\right]
$$

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1 & 1 \\
1 & 0
\end{array}\right]^{2}\left[\begin{array}{l}
F_{n-2} \\
F_{n-3}
\end{array}\right]=\cdots=\left[\begin{array}{ll}
1 & 1 \\
1 & 0
\end{array}\right]^{n}\left[\begin{array}{l}
1 \\
0
\end{array}\right]
$$

We can use SQuaremultPower to compute this!

## Combinatorics

This semester is going to be mainly focused on combinatorics. So let's look at one of the most beautiful combinatorial objects in all of mathematics: Pascal's Triangle

Pascal's Triangle

```
                        1
                        1
        1 2 1
        1 3 3 1
        1 }
        1
        1 
    1
```


## Binomial Coefficients and Pascal's Triangle

- Blaise Pascal first discussed his triangle in his Traité du Triangle Arithmétique [Pas65]
- One of the first works on probability theory


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- Blaise Pascal first discussed his triangle in his Traité du Triangle Arithmétique [Pas65]
- One of the first works on probability theory
- Binomial coefficients were first discussed in detail in India in the tenth-century [Knu97]


## Binomial Coefficients

- "The number of ways to choose $k$ items from $n$ distinct items"

$$
\binom{n}{k}=\frac{n!}{k!(n-k)!}
$$

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- "The number of ways to choose $k$ items from $n$ distinct items"

$$
\binom{n}{k}=\frac{n!}{k!(n-k)!}
$$

- "The number of ways to not choose $n-k$ from $n$ distinct items"

$$
\binom{n}{k}=\binom{n}{n-k}
$$

Pascal's Triangle

$$
\begin{aligned}
& { }^{(0)} \\
& \text { ( }{ }^{1} \text { ) } \quad{ }^{(1)} \\
& \text { (2) (2) (2) } \\
& \text { (3) (3) (3) (3) (3) (3) } \\
& \begin{array}{lllll}
\binom{4}{0} & \binom{4}{1} & \binom{4}{2} & \binom{4}{3} & \binom{4}{0}
\end{array}
\end{aligned}
$$

Pascal's Triangle

$$
\begin{aligned}
& \left({ }_{0}^{0}\right) \\
& \binom{1}{0} \quad\binom{1}{1} \\
& \begin{array}{lll}
\binom{2}{0} & \binom{2}{1} & \binom{2}{2}
\end{array} \\
& \left(\begin{array}{lll}
\left.\begin{array}{l}
3 \\
0
\end{array}\right) & \binom{3}{1} & \binom{3}{2}
\end{array} \quad\binom{3}{3}\right. \\
& \left.\begin{array}{lllllllll}
\binom{4}{0} & \binom{4}{1} & \binom{4}{2} & \binom{4}{3} & \binom{4}{4} & 1 & 4 & 6 & 4
\end{array}\right)
\end{aligned}
$$

## A Pattern in the Triangle

```
                                    1
                                    1
                                    1 2 1
            1 3 3 1
            1 }406\mp@code{6
            1
            1
1
```

A Pattern in the Triangle


## Proving the Pattern

## Claim:

$$
\sum_{k=0}^{\left\lfloor\frac{n}{2}\right\rfloor}\binom{n-k}{k}=F_{n+1}
$$

We are going to prove this by a combinatorial argument

## Staircases

Question: How many ways are there to climb a staircase going one or two steps at a time?


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We can think of this recursively!

## Steps to Compute Steps

- Let the starting step be step 0 . Assuming we are on step $n \geq 2$, how did we get here?


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## Steps to Compute Steps

- Let the starting step be step 0 . Assuming we are on step $n \geq 2$, how did we get here?
- Either we took a single step from step $n-1$
- Or we took two steps from step $n-2$
- Combining the number of ways to get to step $n-1$ with the number of ways to get to step $n-2$ yields the number of ways to get to step $n$


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- $S_{n}=S_{n-1}+S_{n-2}$


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- Combining the number of ways to get to step $n-1$ with the number of ways to get to step $n-2$ yields the number of ways to get to step $n$
- $S_{n}=S_{n-1}+S_{n-2}$
- How many ways are there to get to step 0? Exactly $1\left(S_{0}=1\right)$


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- Let the starting step be step 0 . Assuming we are on step $n \geq 2$, how did we get here?
- Either we took a single step from step $n-1$
- Or we took two steps from step $n-2$
- Combining the number of ways to get to step $n-1$ with the number of ways to get to step $n-2$ yields the number of ways to get to step $n$
- $S_{n}=S_{n-1}+S_{n-2}$
- How many ways are there to get to step 0? Exactly $1\left(S_{0}=1\right)$
- How many ways are there to get to step 1? Exactly $1\left(S_{1}=1\right)$


## Steps to Compute Steps

- Let the starting step be step 0 . Assuming we are on step $n \geq 2$, how did we get here?
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- $S_{n}=S_{n-1}+S_{n-2}$
- How many ways are there to get to step 0? Exactly $1\left(S_{0}=1\right)$
- How many ways are there to get to step 1? Exactly $1\left(S_{1}=1\right)$
- $S_{n}=F_{n+1}$


## Making Choices

- There is another angle to the staircase problem


## Making Choices

- There is another angle to the staircase problem
- We can just choose which steps to take two steps from, and fill the rest with single steps



## Placing Steps

- We have to choose where to place our steps of size 2
- If we have $n$ steps, how many ways can we place $k$ steps of size 2 ?


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$$
\binom{n-k}{k} \text { ways }
$$

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$$

- How many possible values of $k$ are there?

$$
\left\lfloor\frac{n}{2}\right\rfloor
$$

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- If we have $n$ steps, how many ways can we place $k$ steps of size 2 ?

$$
\binom{n-k}{k} \text { ways }
$$

- How many possible values of $k$ are there?

$$
\begin{gathered}
\left\lfloor\frac{n}{2}\right\rfloor \\
\text { Thus, } \sum_{k=0}^{\left\lfloor\frac{n}{2}\right\rfloor}\binom{n-k}{k}=S_{n}=F_{n+1}
\end{gathered}
$$

Questions?
[Combinatorics] has a relation to almost every species of useful knowledge that the mind of man can be employed upon.

- JAMES BERNOULLI, Ars Conjectandi ("The Art of Conjecturing") (1713)


## Bibliography

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Addison-Wesley, Reading, Mass., third edition, 1997.
圊 Blaise Pascal.
Traité du triangle arithmétique, avec quelques autres petits traitez sur la mesme matière. Par Monsieur Pascal.
G. Desprez, 1665.


[^0]:    ${ }^{1}$ Not that one, the other one

