

The PACE 2023 Challenge: Twin-Width

Anakin



Outline

What is Twin-width?

Computing Twin-width

PACE 2023



Section 1

What is Twin-width?



What is Twin-width?

- A large part of combinatorics is defining useful measures of the complexity of structures



What is Twin-width?

- A large part of combinatorics is defining useful measures of the complexity of structures
 - ▶ How easy is it to construct the object?



What is Twin-width?

- A large part of combinatorics is defining useful measures of the complexity of structures
 - ▶ How easy is it to construct the object?
 - ▶ Algorithmic Complexity



What is Twin-width?

- A large part of combinatorics is defining useful measures of the complexity of structures
 - ▶ How easy is it to construct the object?
 - ▶ Algorithmic Complexity
 - ▶ Efficient encodings

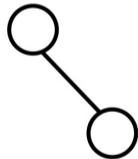
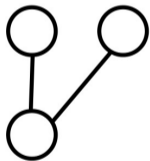


What is Twin-width?

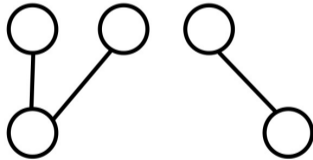
- A large part of combinatorics is defining useful measures of the complexity of structures
 - ▶ How easy is it to construct the object?
 - ▶ Algorithmic Complexity
 - ▶ Efficient encodings
 - ▶ *Decomposition*



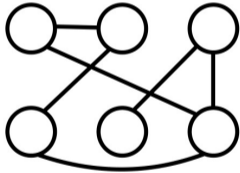
Operations on Graphs: Disjoint Union



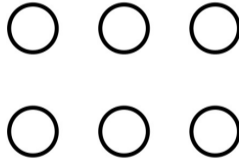
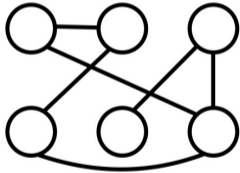
Operations on Graphs: Disjoint Union



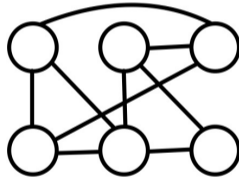
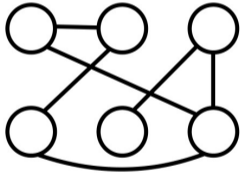
Operations on Graphs: Complement



Operations on Graphs: Complement



Operations on Graphs: Complement



Cographs

We now construct the class of *cographs*



Cographs

We now construct the class of *cographs*

- K_1 is a cograph
- The disjoint union of two cographs is also a cograph
- The complement of a cograph is a cograph



Relating Other Graphs to Cographs

- This efficient way of constructing and decomposing these graphs is useful for many algorithms
 - ▶ Example: Finding the largest complete subgraph of a graph



Relating Other Graphs to Cographs

- This efficient way of constructing and decomposing these graphs is useful for many algorithms
 - ▶ Example: Finding the largest complete subgraph of a graph
- However, not every graph is a cograph
- Can we generalize this by adding some sort of *error* measure?



Relating Other Graphs to Cographs

- This efficient way of constructing and decomposing these graphs is useful for many algorithms
 - ▶ Example: Finding the largest complete subgraph of a graph
- However, not every graph is a cograph
- Can we generalize this by adding some sort of *error* measure?
- This is called *twin-width* [BKTW20]



Contractions

- Throughout this presentation, $G = (V, E)$ will be a connected undirected graph



Contractions

- Throughout this presentation, $G = (V, E)$ will be a connected undirected graph
- Before we can talk about twin-width, we first talk about *contractions* of a graph

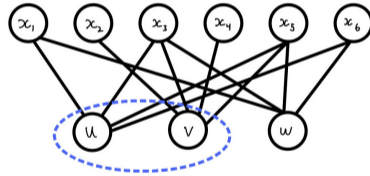
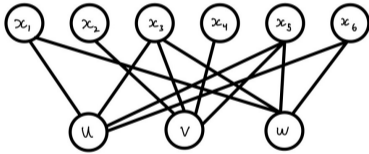


Contractions

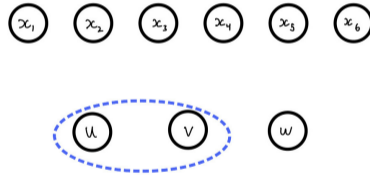
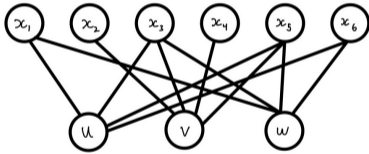
- Throughout this presentation, $G = (V, E)$ will be a connected undirected graph
- Before we can talk about twin-width, we first talk about *contractions* of a graph
 - ▶ Our edges in E will have a color associated with them: black or red.
 - ▶ The red edges will be our *error* that we want
 - ▶ Vertices are black neighbors if linked by a black edge, and red neighbors if linked by a red edge



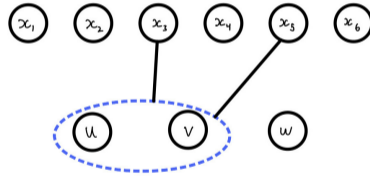
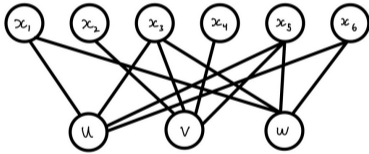
Contractions by picture



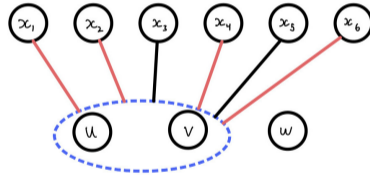
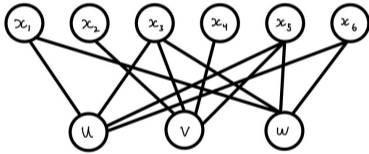
Contractions by picture



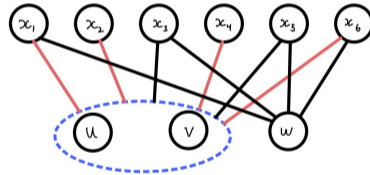
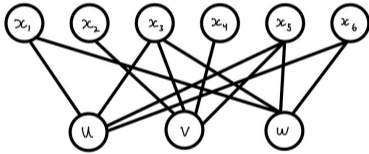
Contractions by picture



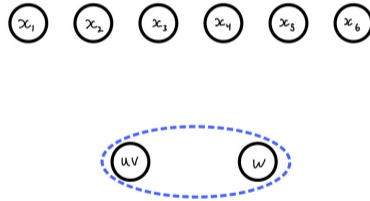
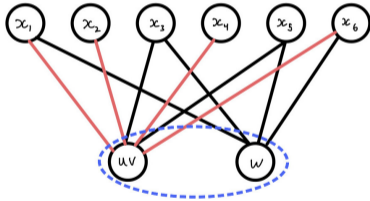
Contractions by picture



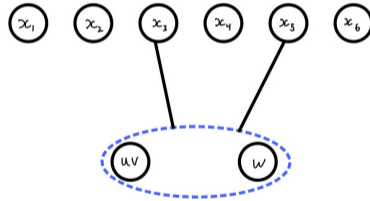
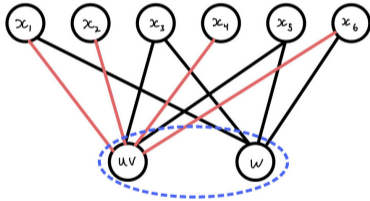
Contractions by picture



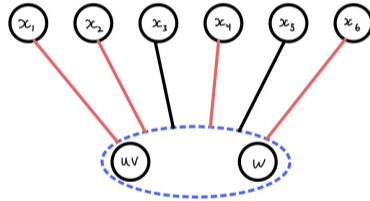
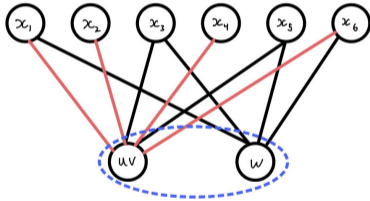
Contractions by picture



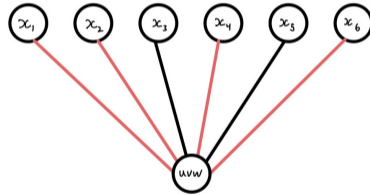
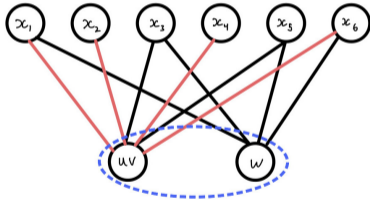
Contractions by picture



Contractions by picture



Contractions by picture



Formal Definition of Contraction

Suppose we are contracting two nodes u, v in G into one node called uv .

- For every neighbor x of u or v :



Formal Definition of Contraction

Suppose we are contracting two nodes u, v in G into one node called uv .

- For every neighbor x of u or v :
 - ▶ If x is a black neighbor of both u and v , then x is now a black neighbor of uv
 - ▶ Otherwise, x is a red neighbor of u and v



Formal Definition of Contraction

Suppose we are contracting two nodes u, v in G into one node called uv .

- For every neighbor x of u or v :
 - ▶ If x is a black neighbor of both u and v , then x is now a black neighbor of uv
 - ▶ Otherwise, x is a red neighbor of u and v
- All other edges are left alone and maintain their color



Twin-width

- Repeatedly applying the contraction operation to nodes of G produces a *contraction sequence* of graphs
 $G = G_n \rightarrow G_{n-1} \rightarrow \cdots \rightarrow G_2 \rightarrow G_1 = K_1$



Twin-width

- Repeatedly applying the contraction operation to nodes of G produces a *contraction sequence* of graphs
 $G = G_n \rightarrow G_{n-1} \rightarrow \cdots \rightarrow G_2 \rightarrow G_1 = K_1$
- The sequence is a *d-sequence* if the maximum number of red edges in any of the G_i in the contraction sequence is d



Twin-width

- Repeatedly applying the contraction operation to nodes of G produces a *contraction sequence* of graphs
 $G = G_n \rightarrow G_{n-1} \rightarrow \cdots \rightarrow G_2 \rightarrow G_1 = K_1$
- The sequence is a *d-sequence* if the maximum number of red edges in any of the G_i in the contraction sequence is d
- The *twin-width* of G is the minimum d such that there exists a d -sequence of G



Observations

- The graph remains connected as we do this



Observations

- The graph remains connected as we do this
- This process will terminate (finitely many nodes and edges)



Observations

- The graph remains connected as we do this
- This process will terminate (finitely many nodes and edges)
- The number of red edges may increase or decrease



Why do We Care?

- We can speed up certain algorithms



Why do We Care?

- We can speed up certain algorithms
- If we have the d -sequence for a graph, we can decompose the graph into complete bipartite graphs and do breadth-first-search in $O(n \log n)$ time [BGK⁺20]
- This works even if the number of edges is $O(n^2)$



Section 2

Computing Twin-width



It's Hard

- There are very few practical algorithms for computing the twin-width of a graph in general



It's Hard

- There are very few practical algorithms for computing the twin-width of a graph in general
- There are some things known for special cases
 - ▶ Cographs are the graphs with twin-width zero
 - ▶ d -dimensional graphs have twin-width $\leq 3d$ [BKTW20]
 - ▶ Planar graphs have twin-width ≤ 9 , and bipartite planar graphs have twin-width ≤ 6 [Hli22]
 - ▶ Graphs with 4 vertices have twin-width ≤ 1 and graphs with 5 vertices have twin-width ≤ 2 [Das22]
 - ▶ In general, bipartite graphs may have arbitrarily large twin-width



SAT Solvers

- Given a boolean formula consisting of variables, **and** gates, and **or** gates, we want to find a satisfying assignment



SAT Solvers

- Given a boolean formula consisting of variables, **and** gates, and **or** gates, we want to find a satisfying assignment
- This is a problem that appears in many places, so readily available *SAT Solvers* exist



SAT Solvers

- Given a boolean formula consisting of variables, **and** gates, and **or** gates, we want to find a satisfying assignment
- This is a problem that appears in many places, so readily available *SAT Solvers* exist
- [SS21] found a way to encode twin-width into a SAT formula
- This is one of the only ways we can reasonably compute twin-widths



Section 3

PACE 2023



A Programming Competition

- Parameterized **A**lgorithms and **C**omputational **E**xperiments is a long-term programming competition
- Our goal will be to devise an efficient algorithm to compute the twin-width d of arbitrary graphs and their d -sequence
 - ▶ Exact Track: Compute the exact sequence
 - ▶ Heuristic Track: Compute an approximate sequence








The Club Submission

- I only found out about this recently so we are a bit behind
- We should make one submission on one track
 - ▶ I was thinking about focusing on the exact track
- There are 100 public test cases + 100 test cases
- Score = the number of test cases we can solve in a given time limit
- My goal is just to put together something and see how well we do



Bibliography

-  Édouard Bonnet, Colin Geniet, Eun Jung Kim, Stéphan Thomassé, and Rémi Watrigant.
Twin-width III: max independent set and coloring.
CoRR, abs/2007.14161, 2020.
-  Édouard Bonnet, Eun Jung Kim, Stéphan Thomassé, and Rémi Watrigant.
Twin-width I: tractable FO model checking.
CoRR, abs/2004.14789, 2020.
-  Kajal Das.
Computation of twin-width of graphs, 2022.
-  Petr Hliněný.
Twin-width of planar graphs is at most 9, and at most 6 when bipartite planar, 2022.
-  André Schidler and Stefan Szeider.
A SAT approach to twin-width.
CoRR, abs/2110.06146, 2021.

